

K22 - Operating Systems: Design Principles and Internals

Fall 2025 @dit

Vaggelis Atlidakis

References: Similar OS courses @Columbia, @Stanford, @UC San Diego, @Brown, @di (previous years);
and textbooks: Operating Systems: Three Easy Pieces, Operating Systems: Principles and Practice, Operating
System Concepts, Linux Kernel Development, Understanding the Linux Kernel

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What you probably ordered?



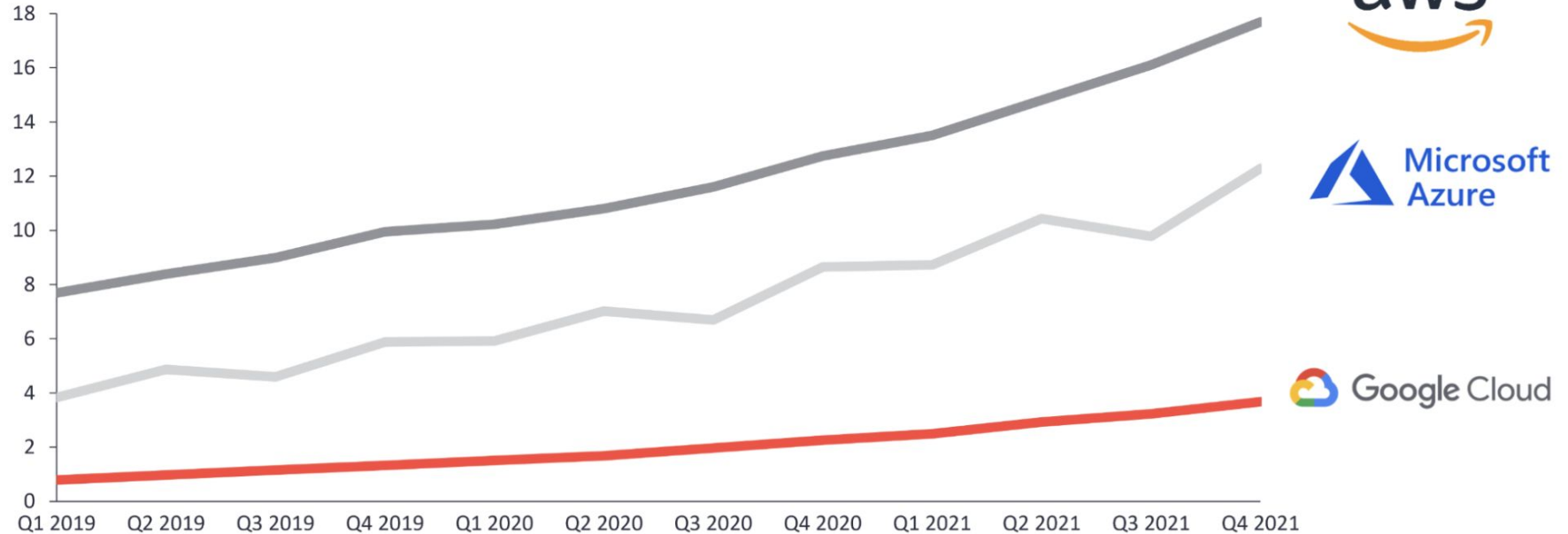
What you ended up getting?



Why study OSes?

Revenue of leading cloud vendors (2019-21)

Quarterly cloud revenue in \$B (IaaS, PaaS, and Others)



[1] <https://iot-analytics.com/cloud-market/>

Why study OSes?

- **These days:** Code developed with superficial understanding
 - No free lunch: Bugs won't fix themselves in production
- **In the near future:** Requirement for software engineers with legit systems understanding will soon be at an all-time high
 - A line of code is not just a line of code
 - Myriads of things happen under the hood

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 - A line of code is not just a line of code
 - Myriads of things happen under the hood
- Learning to navigate a complex codebase, such as an OS kernel, will make you a **fearless** software engineer!

What you need to do to pass the class?

- **Programming assignments (groups of three): 50%**
 - You'll add features to the Linux kernel and run a VM with it
 - You'll need to understand lots of code; won't need to write much
 - **Prerequisites**
 - Do you have a laptop that can run a Linux VM?
 - Can you make teams of three ASAP?
 - Can you use git?
- **Exams: 50%**
 - Midterm: 20%
 - Final: 30%

What you really need to do?

- Come to class, ask questions, be curious
- Choose a team you can collaborate well with
- Own your code
 - You know what you must not do...
 - This is not a class you pass by outsourcing your work
 - There is no magic: You do the work $\Rightarrow$ You pass the class
 - GenAI usage: Use LLMs to help you understand
- It's a difficult course. Hang in there, we will help...

Overview

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 - Intro [Q: What is an OS?]
 - Events [Q: When does the OS run?]
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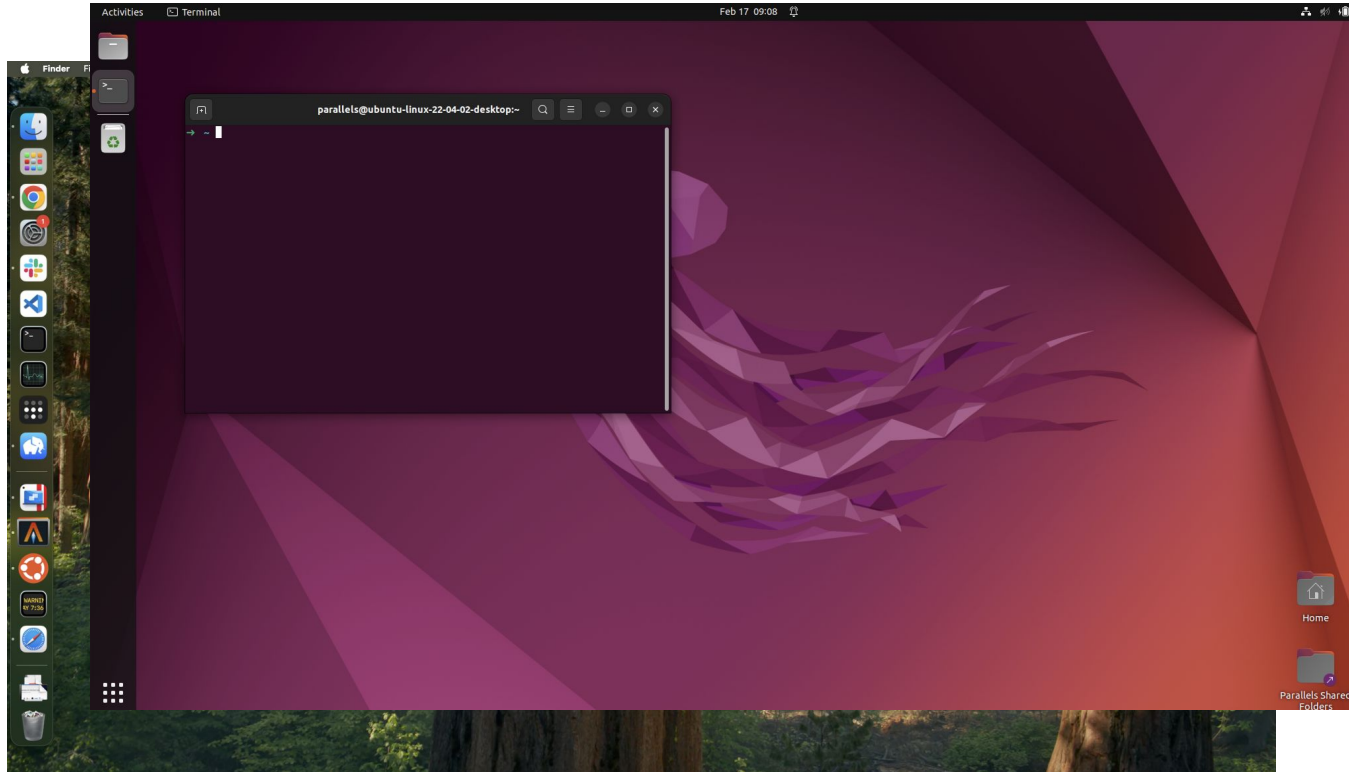
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 - Time Management [Q: What is scheduling?]
 - Memory Management [Q: What is virtual memory?]
 - Files [Q: What is a file descriptor?]
 - Storage Management [Q: How do we allocate disk space to files?]

- * Basic (H/W & S/W)
- * **Abstractions**
- * **Primitives**
- * **Mechanisms**

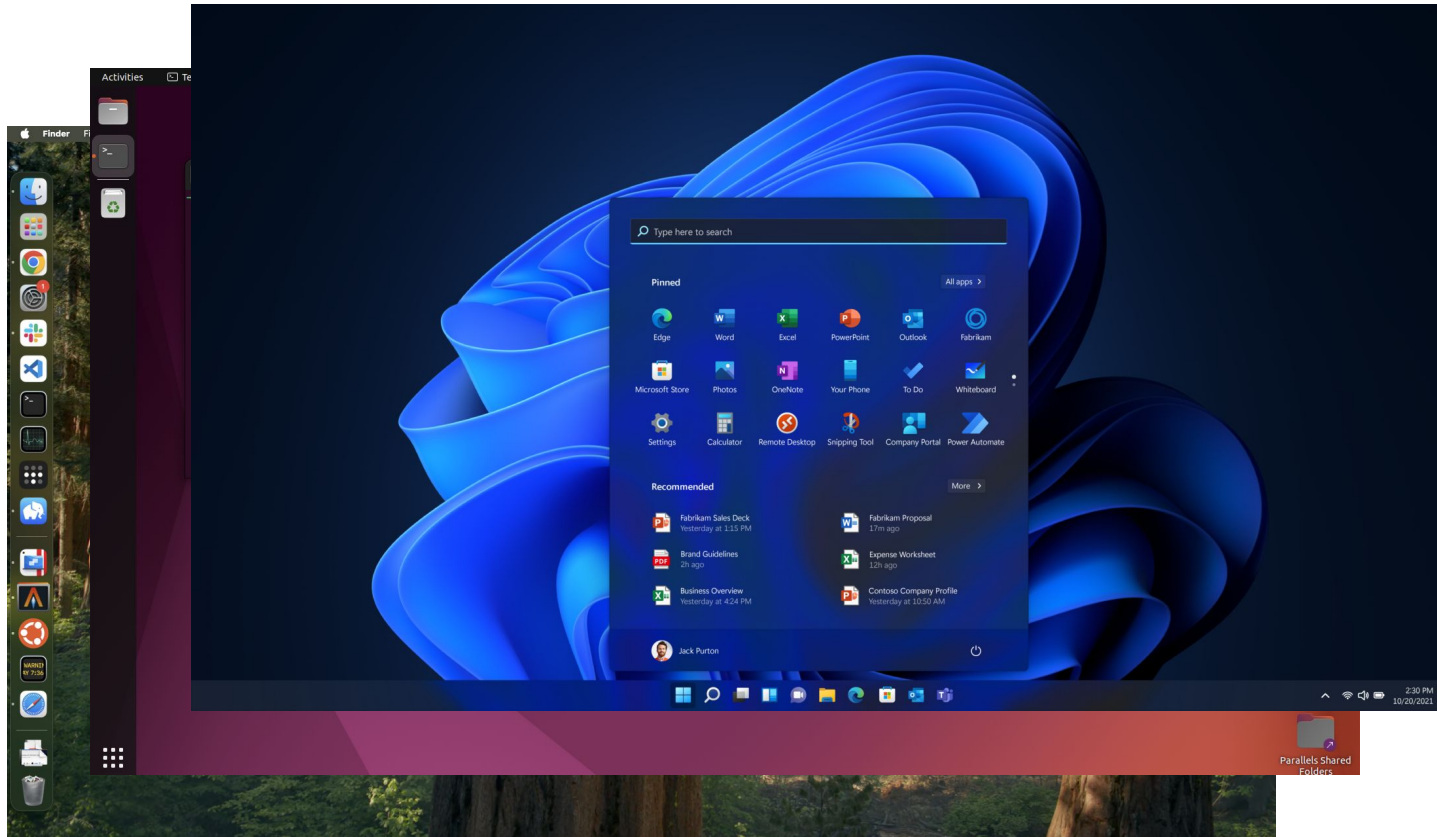
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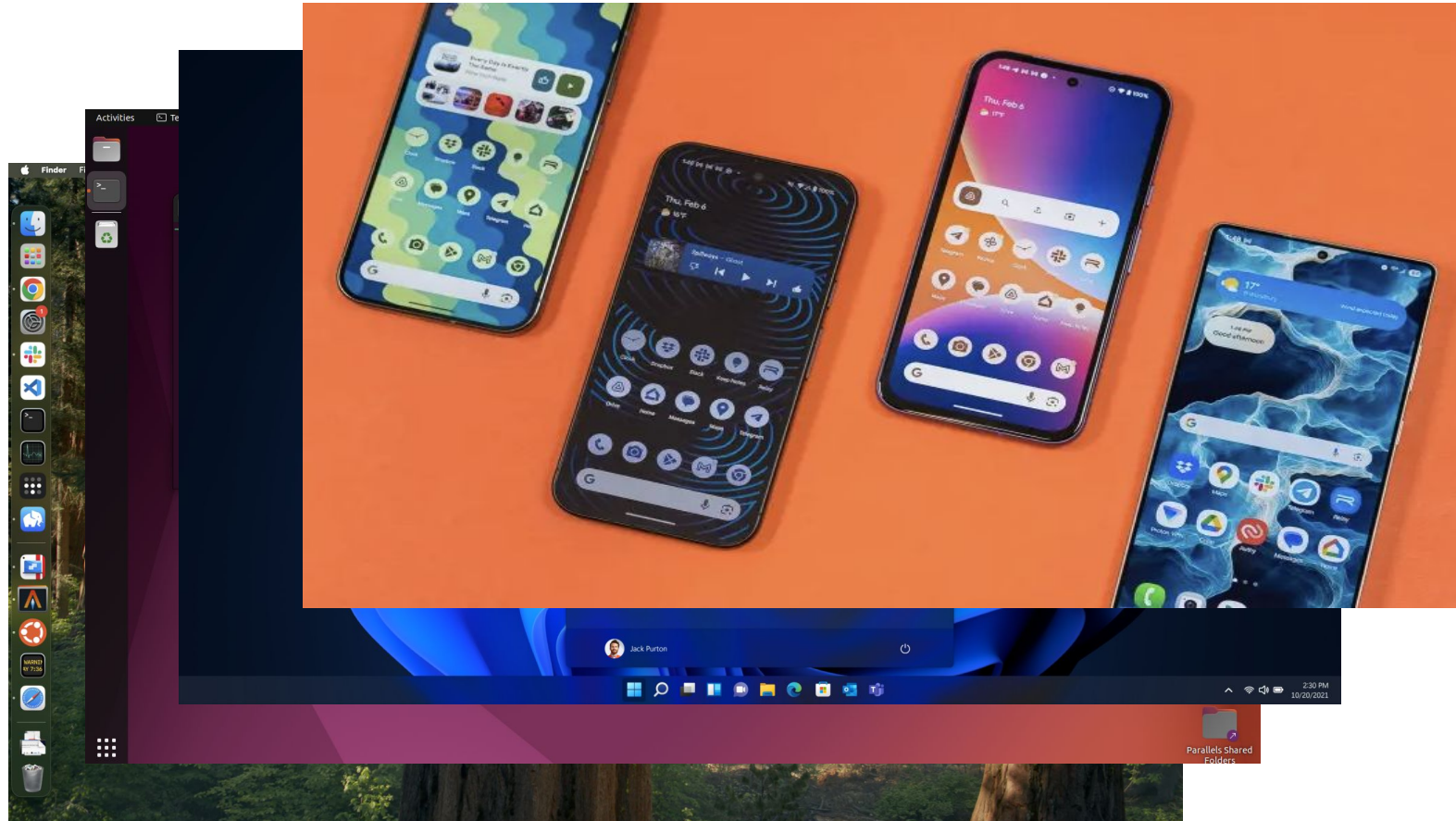
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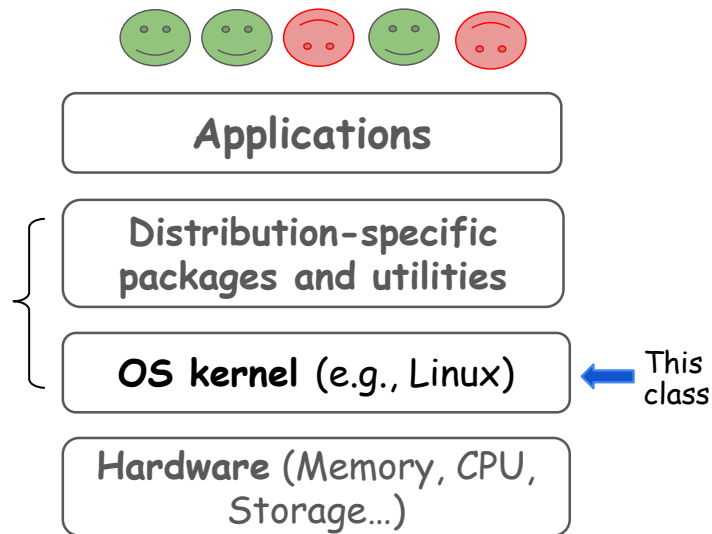
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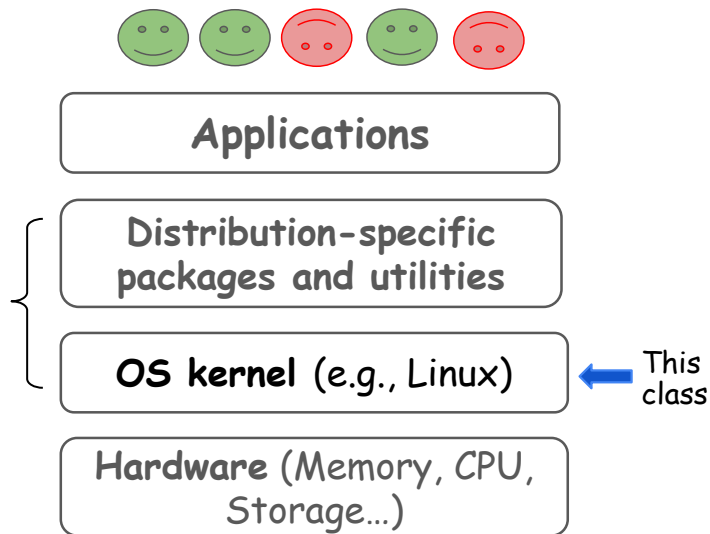


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- Multiple users not necessarily sane, build and run different applications, at the same time on shared hardware

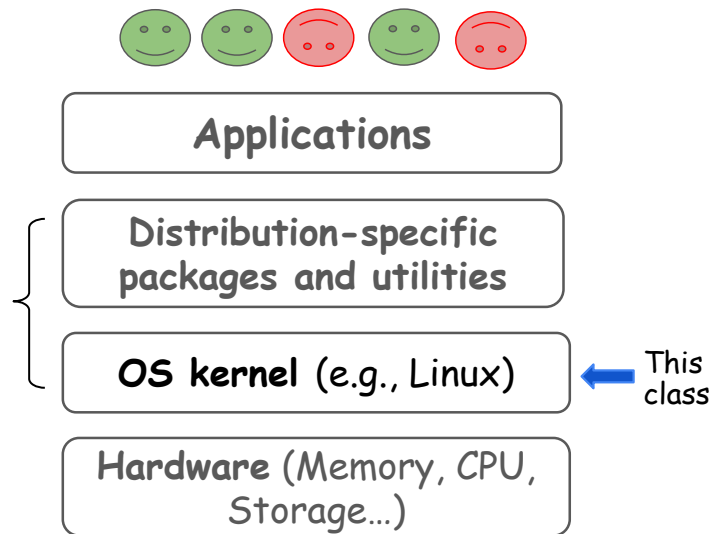


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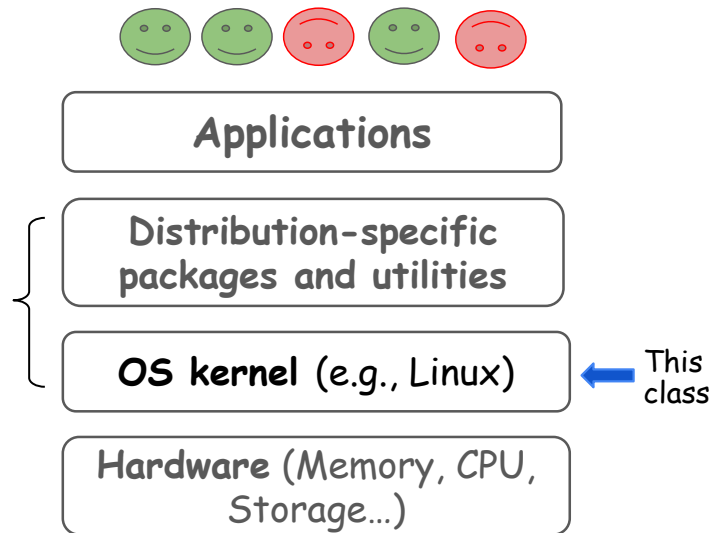
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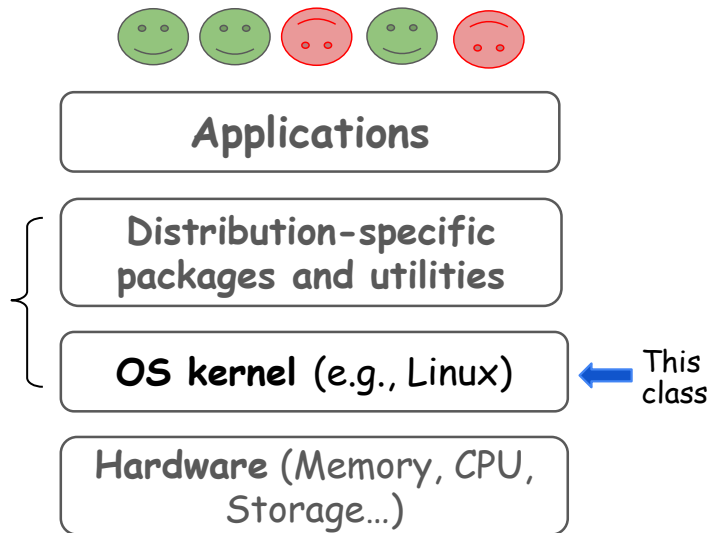
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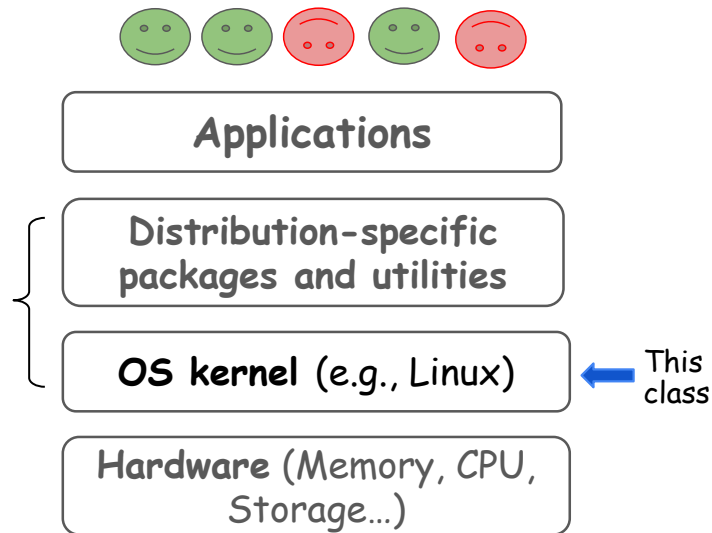
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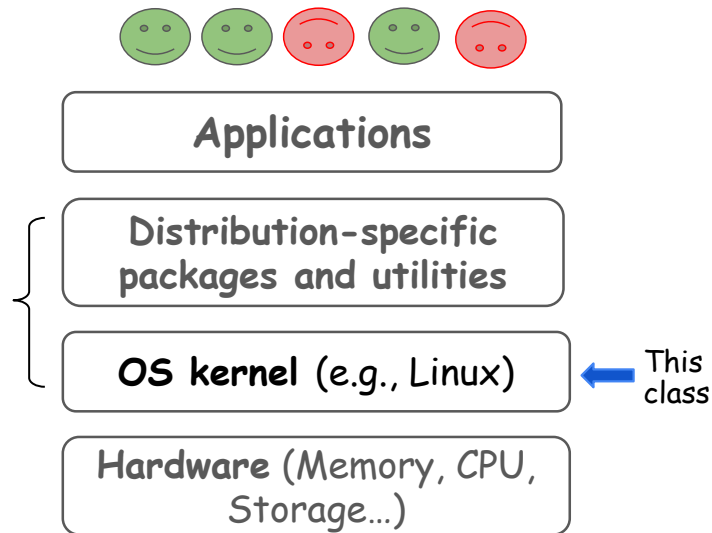
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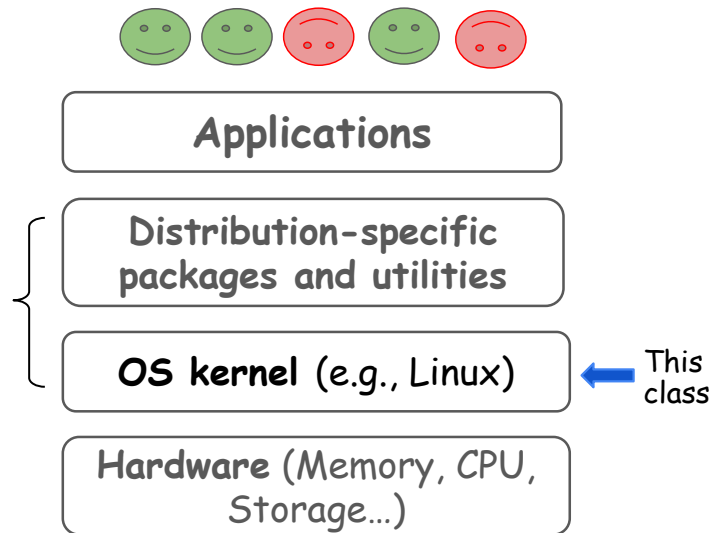
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- **Fairness**: Is the system "pleasant" to use?



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- **Physical Memory (PM)**
 - Stores data addressable on a byte granularity

Physical memory

Addr	Data
0x000	0x1
...	...
0x3FF	0xa
0x400	0x5

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Processor

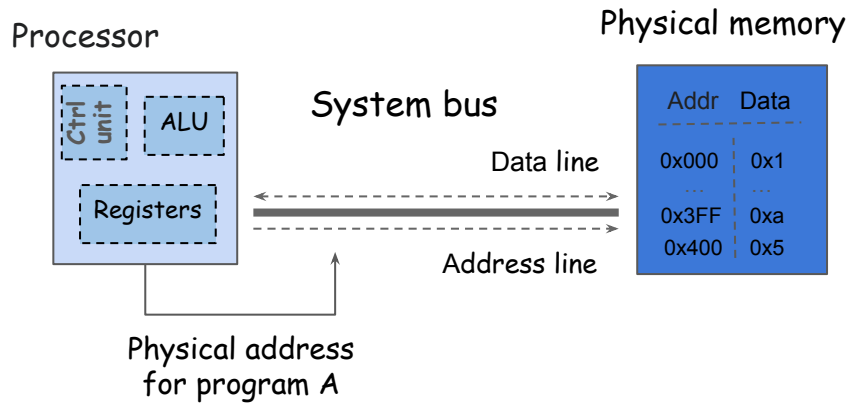


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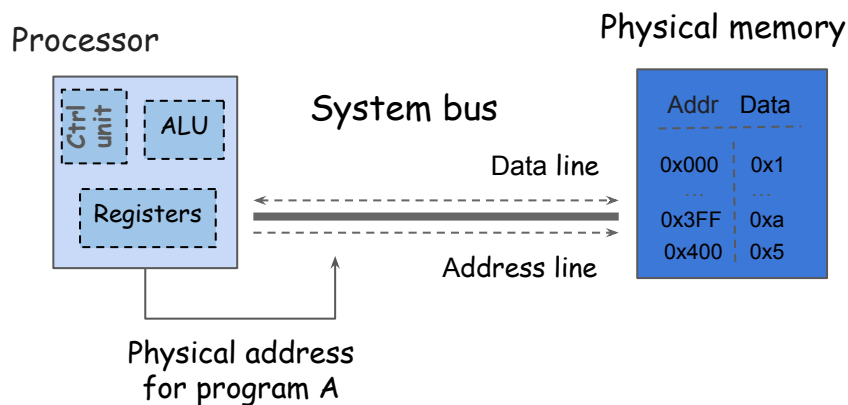
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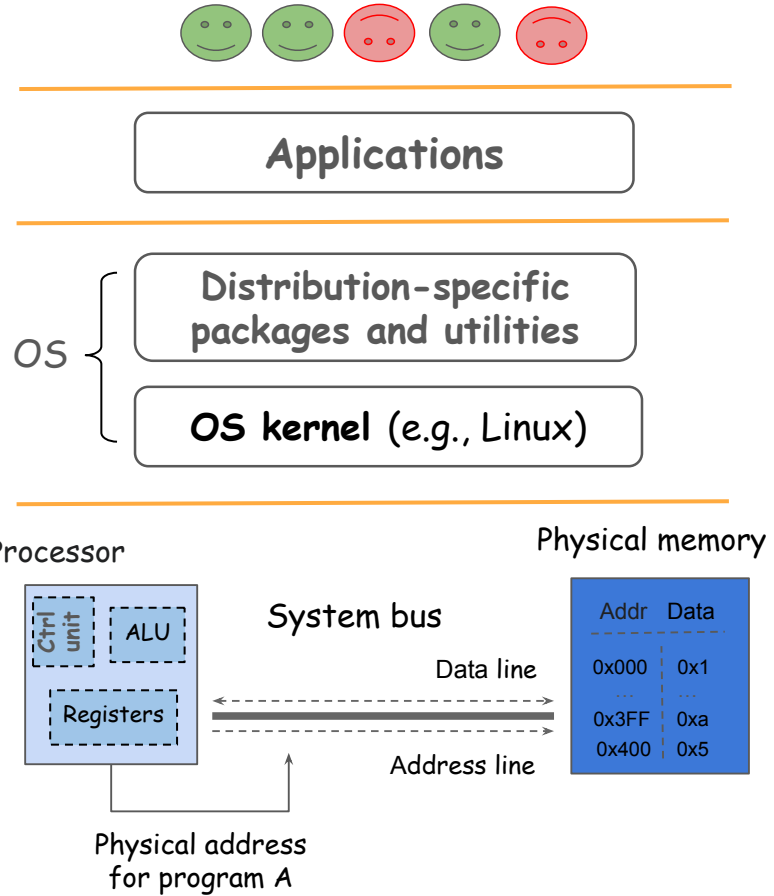


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 - Memory and processor communication channel
- Is this model satisfactory?
- Reminds you of something?

Designing an OS kernel

Desirable properties

1) *Security*, 2) *Reliability*, 3) *Portability*, 4) *Fairness*

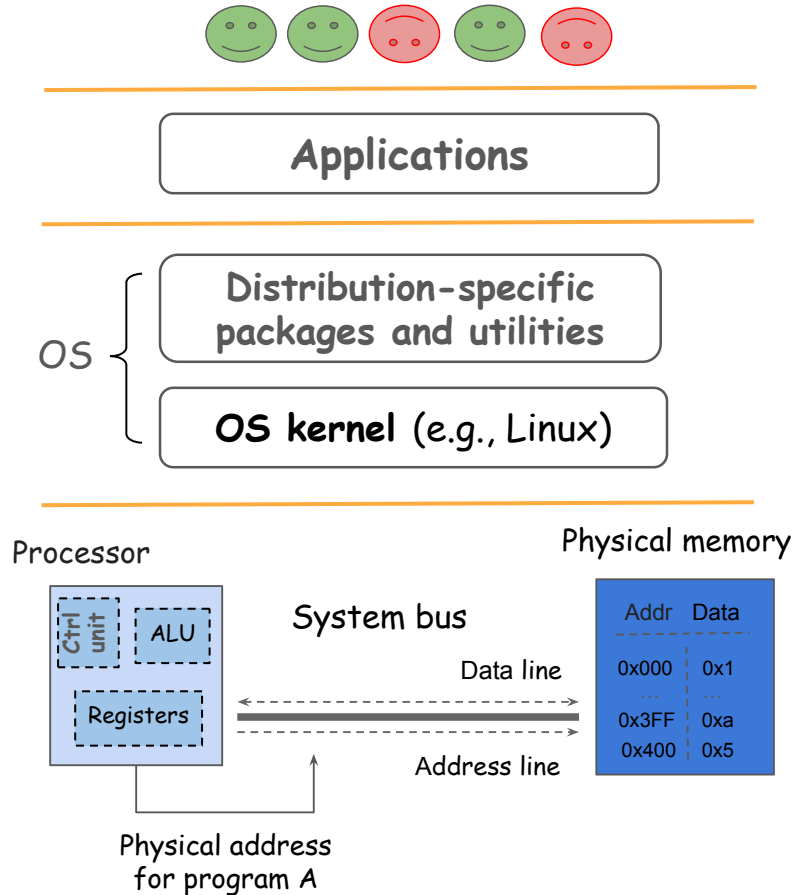


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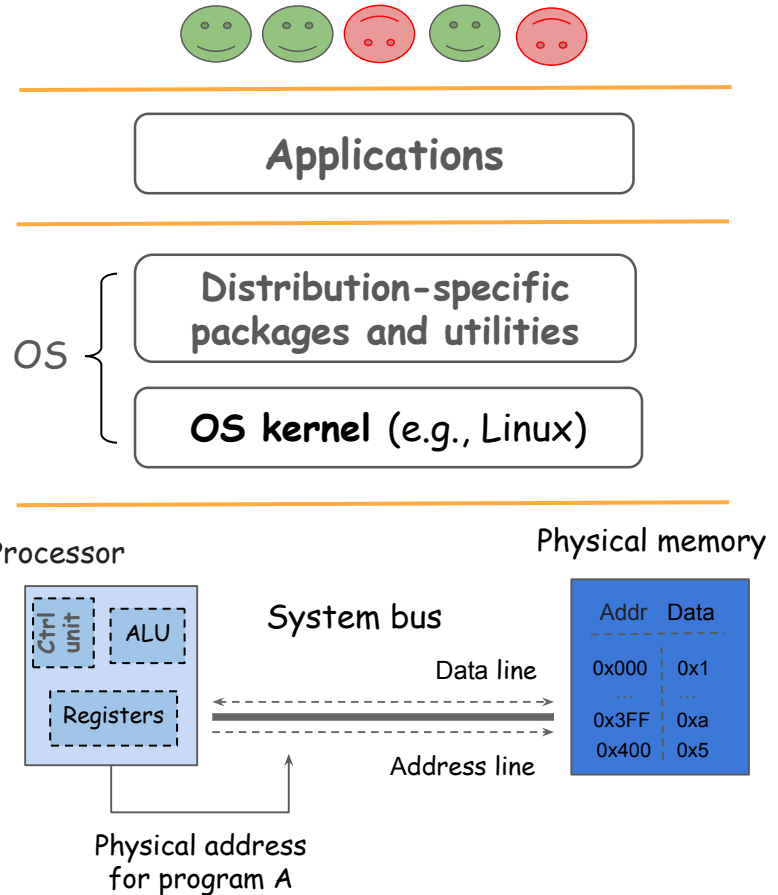
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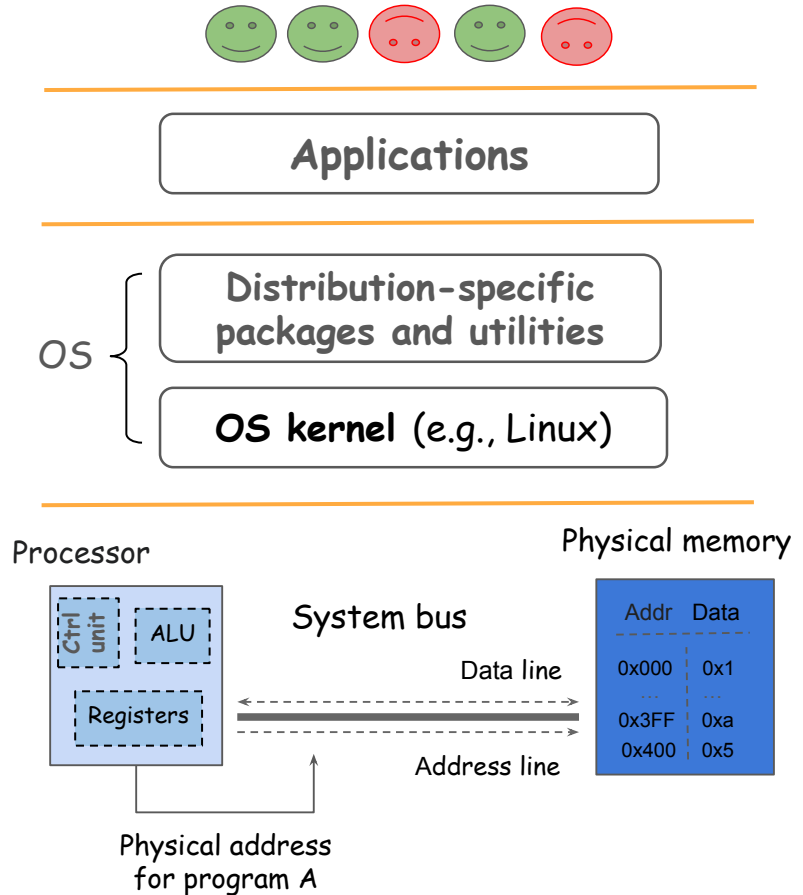
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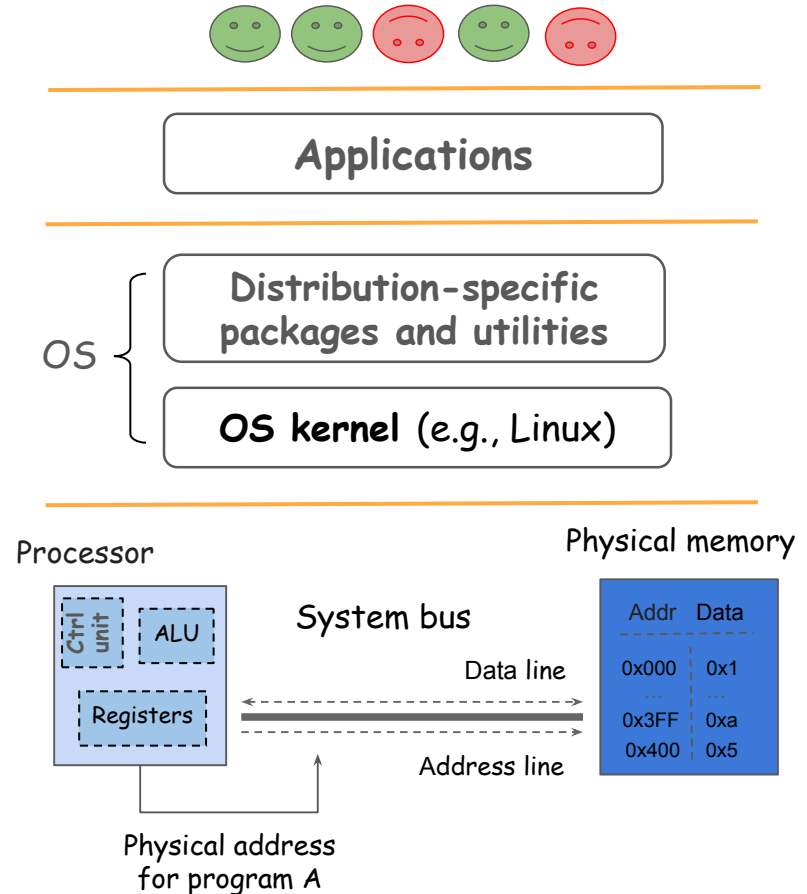
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- **Preemption**: The OS is always able to take control of the processor, regardless of what programs are running

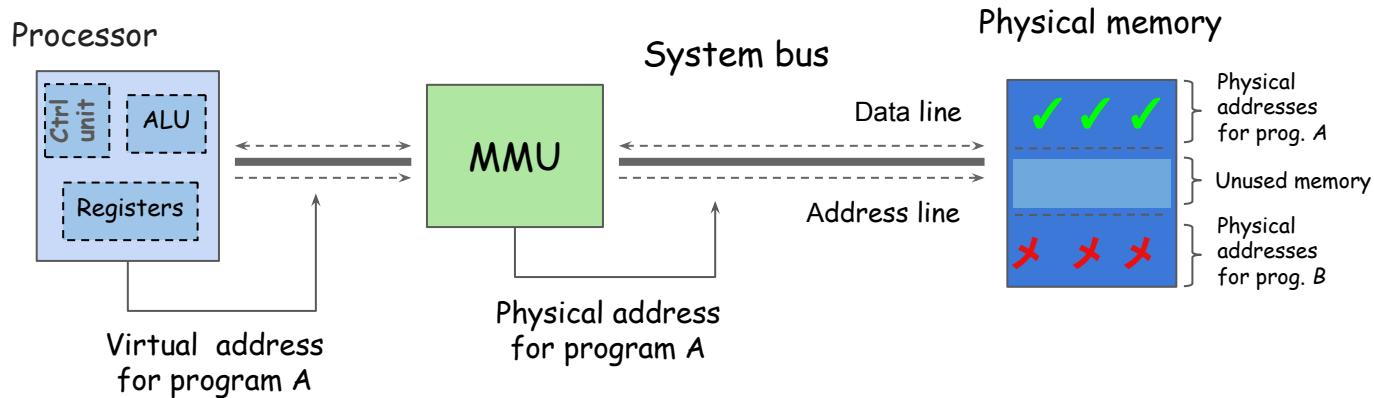


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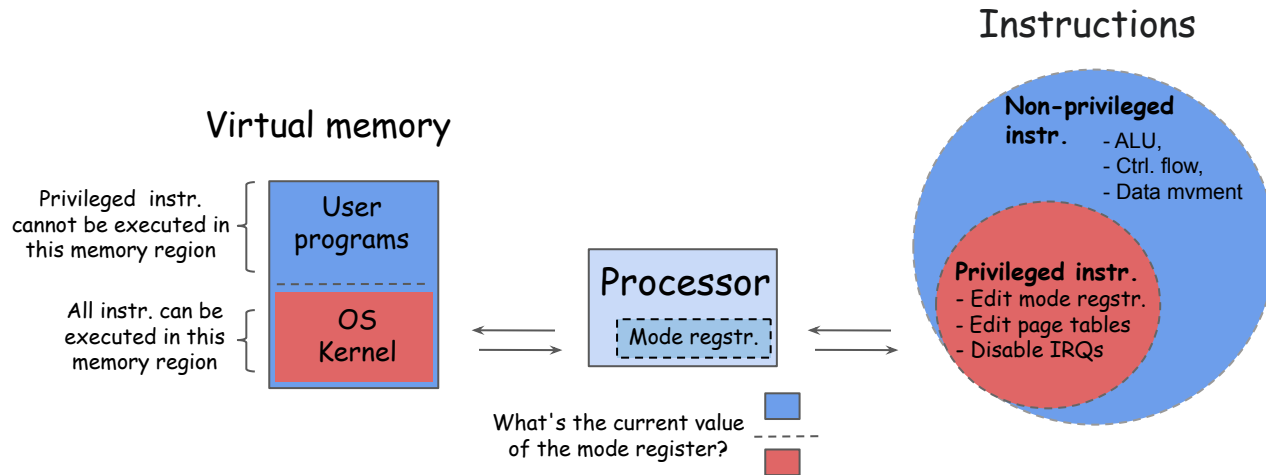


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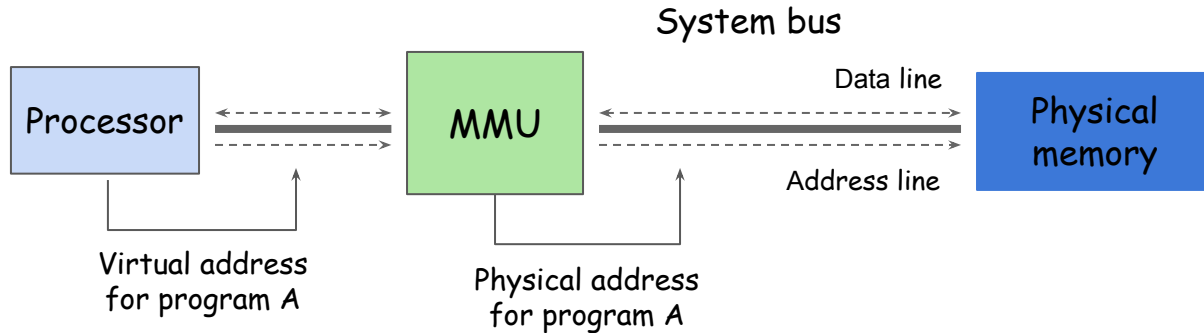


Implementing preemption

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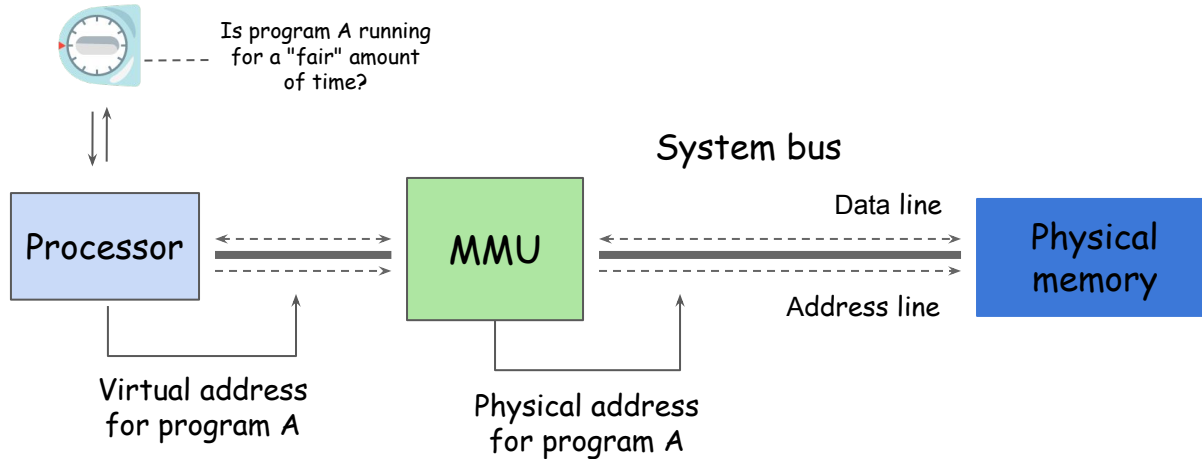
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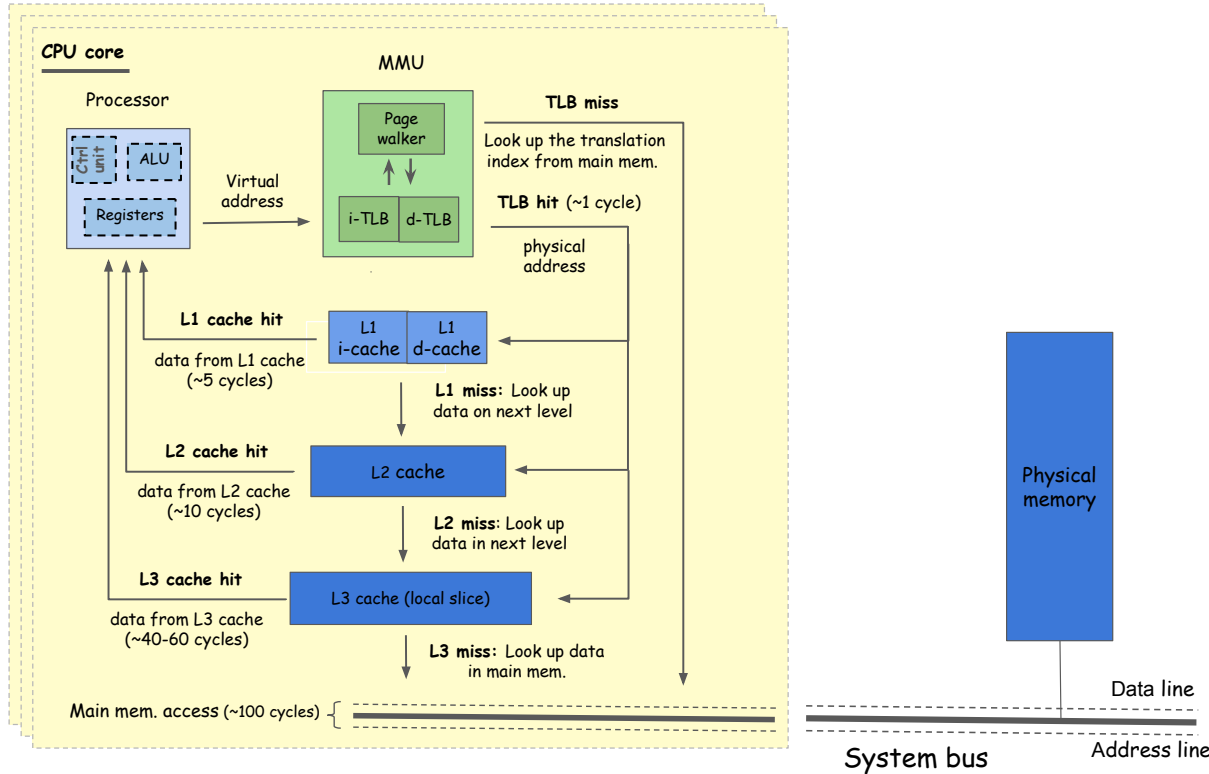


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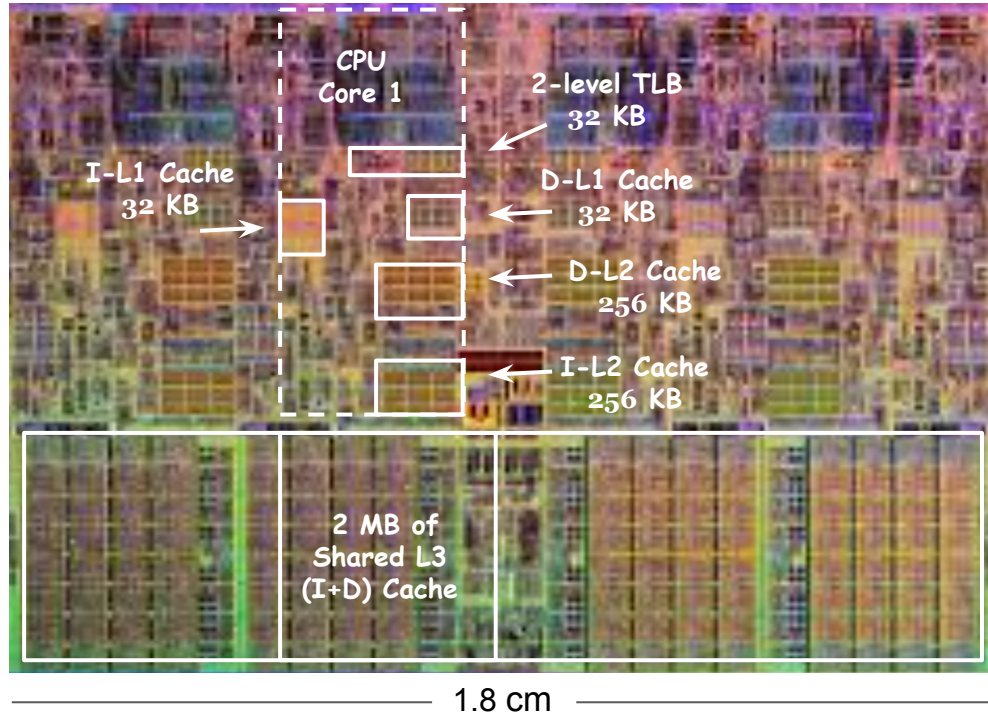


Hardware components leveraged by OS mechanisms



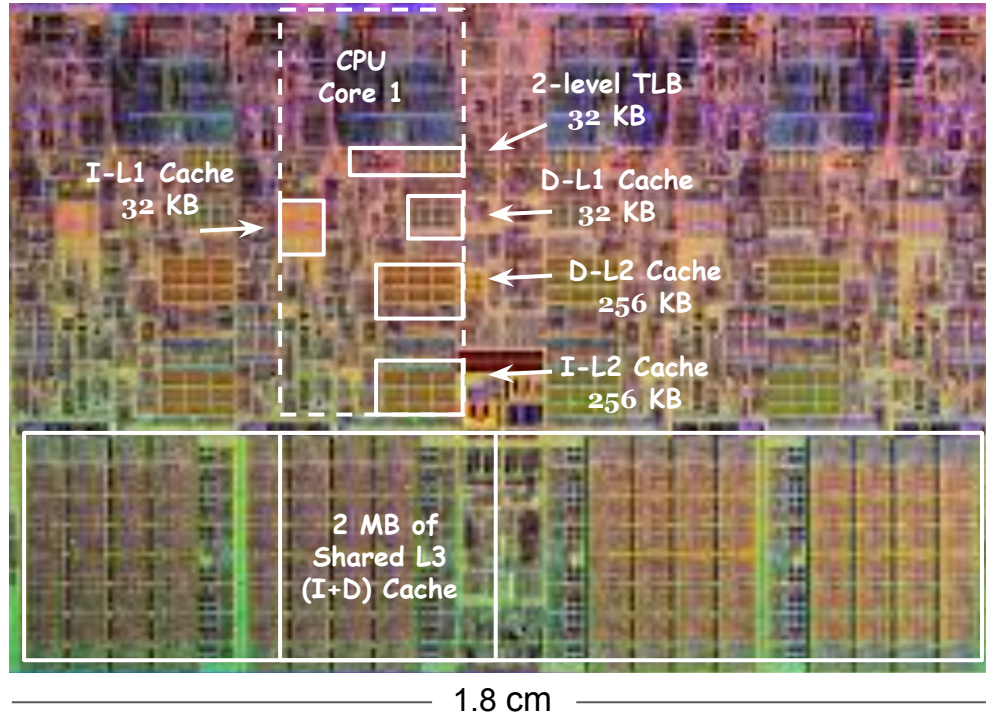
Reality check on hardware's scale

Nahalem Intel® core i7



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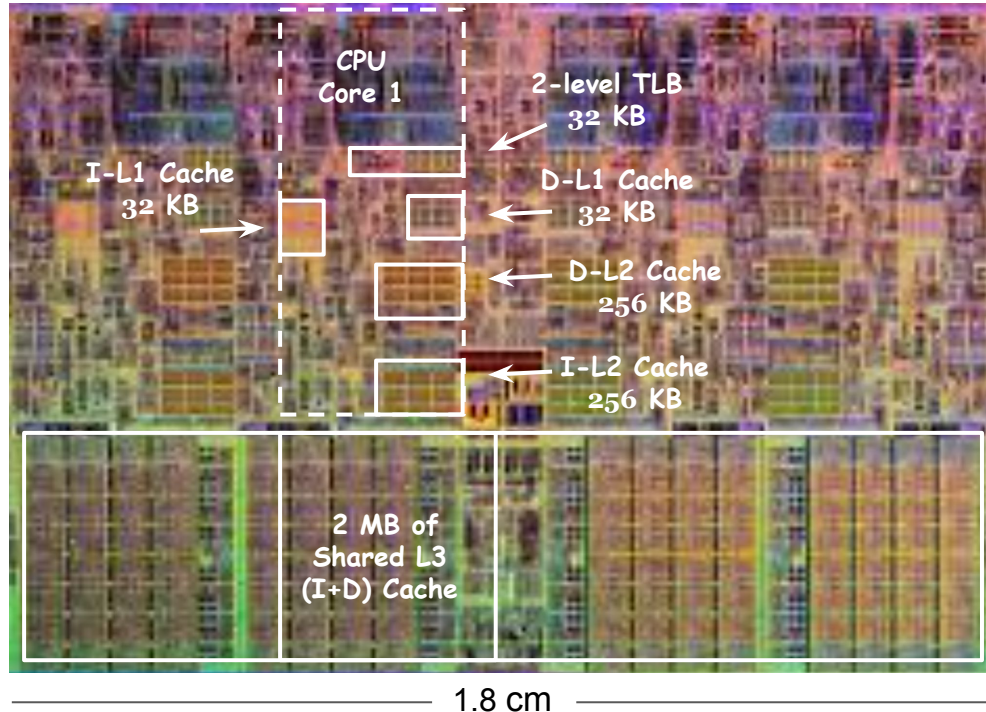
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Access times (1 clock cycle $\sim 0.5\text{ns}$)

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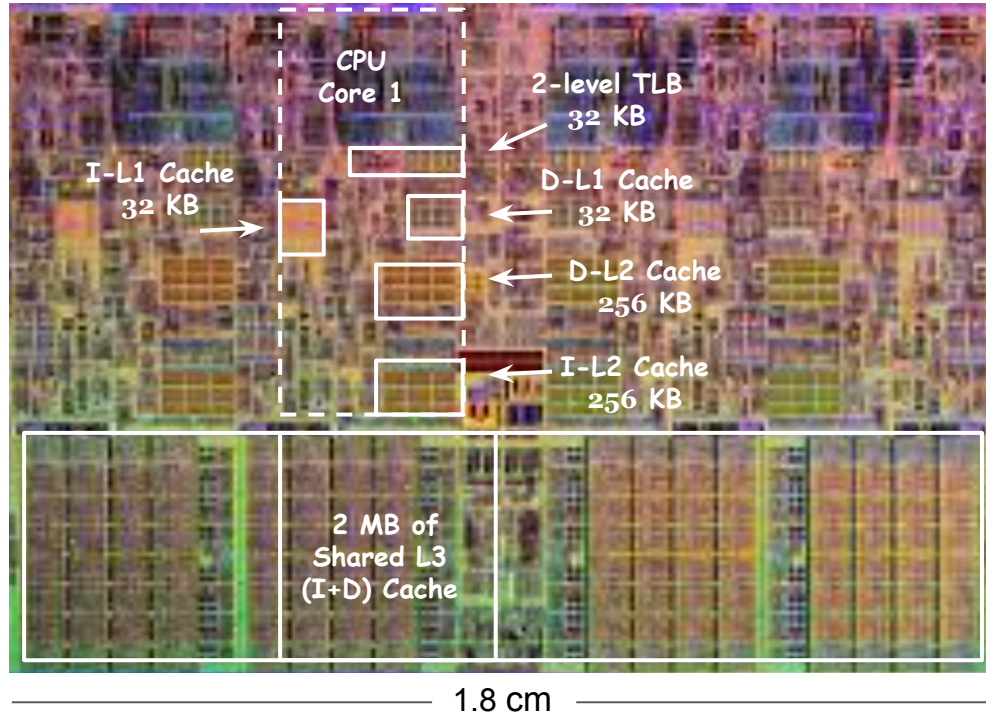


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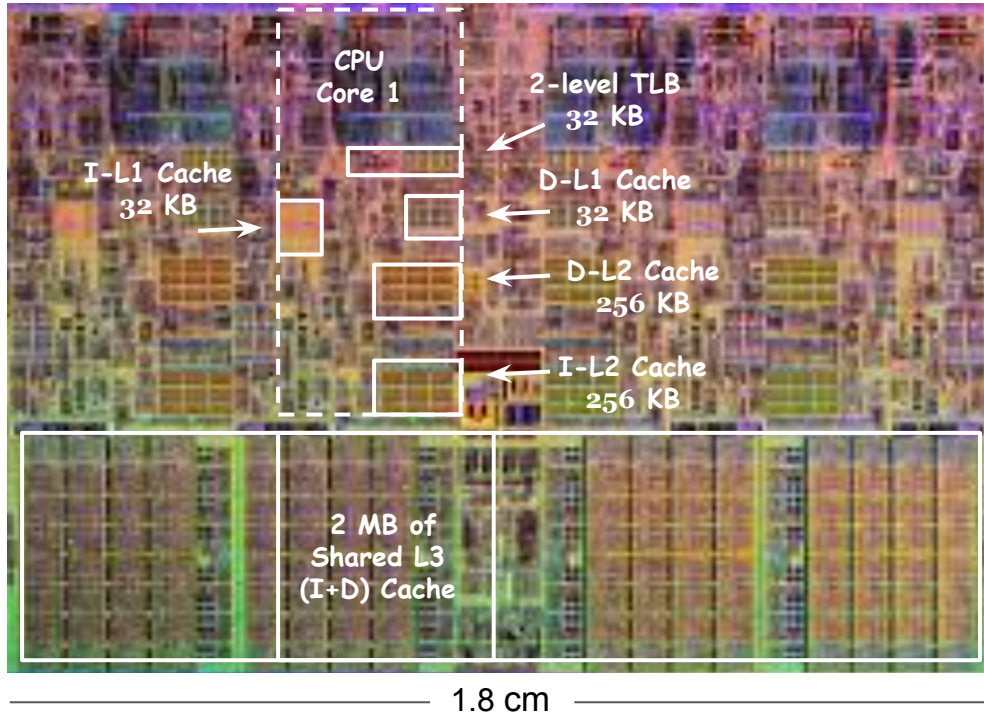


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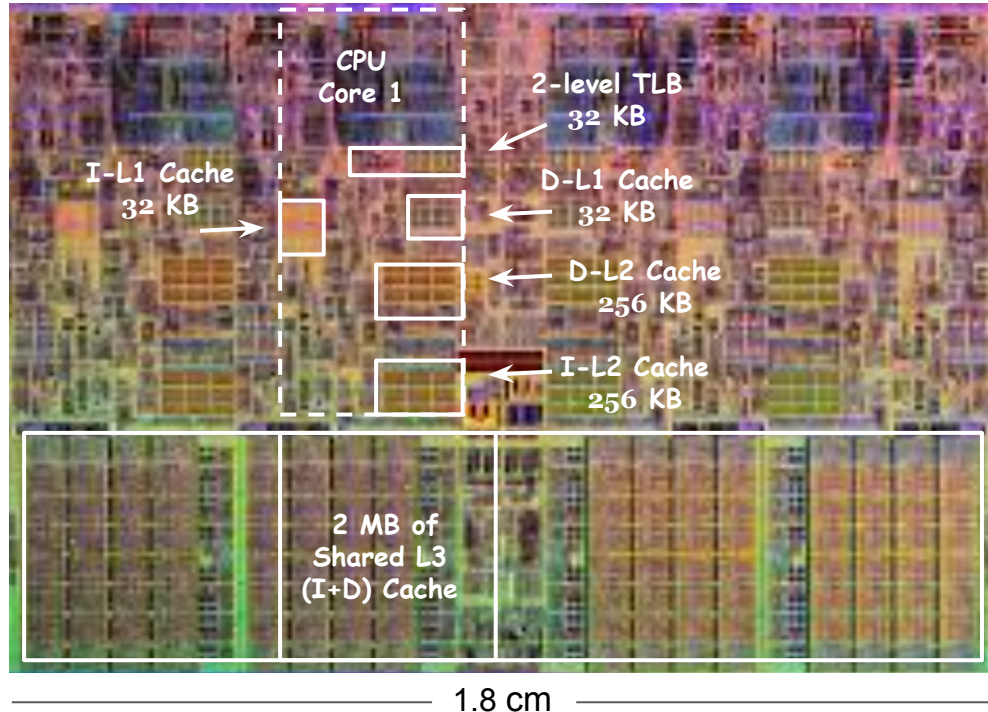


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- L3 Cache miss (MM access): ~100 clock cycles